

$Z_0 \subseteq Z_{\mathcal{L}_{\mathcal{F}}}$ minimal $\mathcal{L}_{\mathcal{F}}$ -theory

containing every thing we need in the following discussion.

$T \subseteq Z_0$ concretely given $\mathcal{L}_{\mathcal{F}}$ -theory
 \downarrow we have the definable set

In T (or in Z_0) $\Gamma T \subseteq \underbrace{W}_{Z_0 \text{ power}} \subseteq V_W$

V_W is a definable class defined by this is a set

$$"x \in V_W" \leftrightarrow (\exists m \in \omega) ("x \in V_m")$$

$\text{Term}_{\mathcal{L}_{\mathcal{F}}}$ the set of all terms | $\text{Func}_{\mathcal{L}_{\mathcal{F}}}$ $\mathcal{L}_{\mathcal{F}}$ -functions
 definable $\mathcal{L}_{\mathcal{F}}$

$\mathcal{C}\text{Term}_{\mathcal{L}_{\mathcal{F}}}$ the definable set of all closed $\mathcal{L}_{\mathcal{F}}$ -terms

Note, for all (concretely given) closed $\mathcal{L}_{\mathcal{F}}$ -term t

$Z_0 \vdash t \in V_W$
 Lemma 17.1 $V_W = \langle V_0, \in, \phi, \{ \cdot, \cdot \}, \dots \rangle$
 In Z_0 we have that letting be the $\mathcal{L}_{\mathcal{F}}$ -structure

$$Z_0 \vdash \forall t \in \mathcal{C}\text{Term}_{\mathcal{L}_{\mathcal{F}}} (t^{V_W} \in V_W)$$

$$Z_0 \vdash (\forall x \in V_W) (\exists t \in \mathcal{C}\text{Term}_{\mathcal{L}_{\mathcal{F}}}) (t^{V_W} \equiv x)$$

For a metamathematical object o we denote with

$\ulcorner o \urcorner_{\mathcal{L}_{\mathcal{F}}}$ term of the corresponding object in set theory

For example, if variables x_0, x_1, \dots are coded as elements of $\omega \times \{0\}$

the $\ulcorner x_0 \urcorner = \{ \{z, z\}, \{z, 0\} \}$

where for $m \in \mathbb{N}$ \underline{m} denotes the numeral for n which is defined inductively as a closed $\mathcal{L}_{\mathcal{F}}$ -term

$$\underline{0} = \phi \quad \underline{m+1} = \underline{m} \cup \{ \underline{m}, \underline{m} \}$$

$\ulcorner \ulcorner \ulcorner \urcorner \urcorner \urcorner$ not an \mathcal{L}_{\exists} -term but rather a definition of a subset of $\mathcal{W}(\mathcal{L}^2)$ (ii) \mathcal{L}_{\exists}

Let $Z_0 \subseteq \mathcal{T}$,
 for $x \in V_{\mathcal{W}}$ let $\ulcorner x \urcorner$ be the standard element of $\mathcal{A}(\text{Term}_{\mathcal{L}_{\exists}})$ s.t.
 $\ulcorner x \urcorner \ulcorner \ulcorner \urcorner = x$

Subst: $\text{Fml}_{\mathcal{L}_{\exists}} \times \mathcal{W} \times \text{Term}_{\mathcal{L}_{\exists}} \rightarrow \text{Fml}_{\mathcal{L}_{\exists}}$

the st.
 $\text{Subst}(\varphi, m, t)$ = the formula obtained from φ by replacing all free appearances of $x_m (= \langle m, s \rangle)$ with the term t

Lemma 7.2
 For any concretely given \mathcal{L}_{\exists} -formula φ and sequence t of symbols and number k

note that $\ulcorner t \urcorner$, k are \mathcal{L}_{\exists} -terms, we have

- (1) $Z_0 \vdash \ulcorner \varphi(\ulcorner t \urcorner / x_m) \urcorner \equiv \text{Subst}(\ulcorner \varphi \urcorner, m, \ulcorner t \urcorner)$
- (2) $Z_0 \vdash \ulcorner \varphi(\ulcorner k \urcorner / x_m) \urcorner \equiv \text{Subst}(\ulcorner \varphi \urcorner, m, \ulcorner k \urcorner)$

Theorem 15.2 (Diagonal Lemma)

For any \mathcal{L}_{\exists} -formula φ (with $x_0 \in \text{free}(\varphi)$) there is an \mathcal{L}_{\exists} -formula σ s.t.
 (1) $\text{free}(\sigma) \subseteq \text{free}(\varphi) \setminus \{x_0\}$

(2) $Z_0 \vdash \sigma \leftrightarrow \varphi(\ulcorner \sigma \urcorner / x_0)$
Proof If $x_0 \notin \text{free}(\varphi)$ then $\sigma = \varphi$ will do.
 So assume $x_0 \in \text{free}(\varphi)$.

Actually we don't need this restriction!

In \mathcal{Z}_0 , let $f^* : (w)(w^2) \rightarrow w)(w^2)$ be defined

by

$$(f) f^*(p, t) = \begin{cases} u & \text{if } p \in \text{Fml } \mathcal{L} \text{ and} \\ & u \equiv \text{Subst}(p, \sigma, t) \end{cases}$$

\emptyset otherwise

Let k be the first index s.t. x_k does not appear in ψ
 Let σ^* be the (concretely given) σ (class)

$$(*) \quad \forall x_k (f^*(x_0, x_0) \equiv x_k \rightarrow \psi(x_k/x_0))$$

Let σ be the \mathcal{L}_{\exists} -sentence

$$(**) \quad \forall x_k (f^*(\sigma^*, \sigma^*) \equiv x_k \rightarrow \psi(x_k/x_0))$$

We have, by (f), we show that σ is a desired.

$$(f) \quad \mathcal{Z}_0 \vdash f^*(\sigma^*, \sigma^*) \equiv \underbrace{\forall x_k (f^*(\sigma^*, \sigma^*) \equiv x_k \rightarrow \psi(x_k/x_0))}_{= [\sigma]} \text{ by } (**)$$

Thus

$$\mathcal{Z}_0 \vdash (\sigma \rightarrow \psi(\sigma/x_0))$$

On the other hand, noting that

$$(\psi(\sigma/x_0) \rightarrow f^*(\sigma^*, \sigma^*) \equiv x_k \rightarrow \psi(\sigma/x_0))$$

is a tautology $(A \rightarrow (B \rightarrow A))$

Since x_k does not appear in ψ

$$\mathcal{Z}_0 \vdash (\psi(\sigma/x_0) \rightarrow \forall x_k (f^*(\sigma^*, \sigma^*) \equiv x_k \rightarrow \psi(\sigma/x_0)))$$

By (f) it follows that

$$\mathcal{Z}_0 \vdash (\psi(\sigma/x_0) \rightarrow \underbrace{\forall x_k (f^*(\sigma^*, \sigma^*) \equiv x_k \rightarrow \psi(x_k/x_0))}_{\sigma})$$

Thus

$$\mathcal{Z}_0 \vdash (\psi(\sigma/x_0) \rightarrow \sigma) \quad \square$$